INTRODUCTION TO CODE ANALYSIS AND OPTIMIZATION

PROGRAM ANALYSIS AND OPTIMIZATION – DCC888

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The goal of a compiler writer is to bridge the gap between programming languages and the hardware; hence, making programmers more productive.

A compiler writer builds bridges between people and machines, and this task is each day more challenging.
"We do not need that many compiler guys. But those that we need, we need them badly."

François Bodin – CEO of CAPS

A lot of amazing people in computer science were working with compilers and programming languages.

Who do you know in these photos?

And many of the amazing things that we have today only exist because of compilers and programming languages!
The primary goal of this course is to explain the student how to transform a program automatically, while preserving its semantics, in such a way that the new program is more efficient according to a well-defined metric.

- There are many ways to compare the performance of programs:
  - Time
  - Space
  - Energy consumption

1) Is a smaller program always faster?
2) Is a faster program always more energy efficient?
Goals of this Course

THE SECOND GOAL OF THIS COURSE IS TO INTRODUCE STUDENTS TO TECHNIQUES THAT LET THEM UNDERSTAND A PROGRAM TO A LEVEL THAT WOULD BE HARDLY POSSIBLE WITHOUT THE HELP OF A MACHINE.

• We understand programs via static analysis techniques.
• These analyses are key to enable code optimizations.
• But they also have other uses:
  – They help us to prove correctness properties.
  – They help us to find bugs in programs.
Goal 1: Program Optimization

- The main goal of the techniques that we will see in this course is to optimize programs.
- There are many, really many, different ways to optimize programs. We will see some of these techniques:
  - Copy elimination
  - Constant propagation
  - Lazy Code Motion
  - Register Allocation
  - Loop Unrolling
  - Value Numbering
  - Strength Reduction
  - Etc, etc, etc.

```c
#include <stdio.h>
#define CUBE(x) (x)*(x)*(x)
int main() {
    int i = 0;
    int x = 2;
    int sum = 0;
    while (i++ < 100) {
        sum += CUBE(x);
    }
    printf("%d\n", sum);
}
```

How can you optimize this program?

```assembly
_main:
    pushl %ebp
    movl %esp, %ebp
    pushl %ebx
    subl $20, %esp
    call L3
L3:
    popl %ebx
    movl $800, 4(%esp)
    movl %eax, (%esp)
    call _printf
    addl $20, %esp
    popl %ebx
    leave
    ret
```
Goal 2: Bug Finding

• Compiler analyses are very useful to find (and sometimes fix) bugs in programs.

```c
void read_matrix(int* data,
                 char w, char h) {
    char buf_size = w * h;
    if (buf_size < BUF_SIZE) {
        int c0, c1;
        int buf[BUF_SIZE];
        for (c0 = 0; c0 < h; c0++) {
            for (c1 = 0; c1 < w; c1++) {
                int index = c0 * w + c1;
                buf[index] = data[index];
            }
        }
    }
    process(buf);
}
```

– Null pointer dereference
– Array out-of-bounds access
– Invalid Class Cast
– Tainted Flow Vulnerabilities
– Integer Overflows
– Information leaks

Can you spot a security bug in this program. **Be aware:** the bug is tricky.
The Contents of the Course

• All the material related to this course is available on-line, at http://www.dcc.ufmg.br/~fernando/classes/dcc888.

• This material includes:
  – Slides
  – Project assignment
  – Class Exercises
  – Useful links

• The page also contains the course bibliography.

Code Analysis and Optimization - DCC888

The goal of this class is to introduce the student to the most recent techniques that compilers use to analyze and optimize programs. The student will learn about dataflow and constraint based program analyses. He or she will have contact with type systems, and the many variants of inductive techniques to prove properties about programs. The class contains a number of expository lectures, and some paper discussion. During the discussions, the students will have the opportunity to get in touch with state-of-the-art techniques published in top conferences in the field of programming languages. The class also features a project assignment, which consists in the implementation of partial redundancy elimination, a classic compiler optimization, in the llvm compiler.

• The course bibliography.
• The syllabus adopted in this course.
• The grading policy.
• The Project Assignment.
• Consulte sua nota computada até agora.

Code: DCC 888
Department: Computer Science
Term: 2013/1
Time: Monday/Wednesday, 5:00pm-6:40pm
Room number: ICEx 2014
Discussion list: dcc888 at googlegroups ...
The Contents of the Course

- The course has 20 lectures. Slides for each lecture are available on-line:

**Lectures**

1. Introduction
2. Loop Optimizations
3. Static Single Assignment
4. Sparse Analyses
5. Range Analysis
6. SSA-Based Register Allocation
7. Just-in-time compilers

**Labs**

1. Intro to LLVM
2. Compiling a language
3. Writing a pass
4. Testing
5. Using the API
6. Interprocedural Analyses
Where we will be

• A compiler has three main parts
  – The front-end is where parsing takes place.
  – The middle-end is where optimizations and analyses take place.
  – The back-end is where actual machine code is generated.

We shall stay here!

In which CS courses can we learn about the other phases?
There is no Silver Bullet

1. It is impossible to build the perfect optimizing compiler.
   • The perfect optimizing compiler transforms each program P in a program $P_{opt}$ that is the smallest program with the same input/output behavior as P.
There is no Silver Bullet

1. It is impossible to build the perfect optimizing compiler.

Let's assume that we are optimizing for size. We can reduce the problem of building the perfect compiler to an undecidable problem, such as the OUTPUT PROBLEM, e.g., "Does the program P output anything?"

The smallest program that does nothing and does not terminate is $P_{\text{least}} = \text{\# L: goto L;}$. By definition, the perfect compiler, when fed with a program that does not generate output, and does not terminate, must produce $P_{\text{least}}$.

Thus, we have a decision procedure to solve the OUTPUT PROBLEM: given a program P, if the perfect compiler transforms it into $P_{\text{least}}$, then the answer to the problem is No, otherwise it is Yes.
The Full-Employment Theorem

1. If C is an optimizing compiler for a *Turing Complete Language*, then it is possible to build a better optimizing compiler than C.

   • In other words, compiler writers will always have a job to do, as it is always possible to improve any existing compiler that compiles any meaningful programming language.

   How could we prove this new theorem?
The Full-Employment Theorem

1. If C is an optimizing compiler for a *Turing Complete Language*, then it is possible to build a better optimizing compiler than C.

   Lets assume that there exists the best compiler B. If P is a program, then let B(P) be its optimized version. There must be a program \( P_x \) that does not terminate, such that \( B(P_x) \neq [L: \text{goto } L] \), otherwise B would be the perfect compiler. As we have seen, this is impossible.

   Therefore, there exists a compiler B' that is better than B, as we can define it in the following way:

   \[
   B'(P) = \begin{cases} 
   [L: \text{goto } L] & \text{if } P = P_x \\
   B(P) & \text{else}
   \end{cases}
   \]
WHY TO LEARN COMPILERS?
The Importance of Compilers

Robert Hundt is a leading compiler engineer working at Google. Read below what he says about the importance of compilers in general, and in that company in particular.

• "At the scale of datacenters, every single performance percent matters! Just take a look at Google's (and other's) publicly available numbers on expenditures on datacenters. We are talking about billions of dollars. A single percent improvement can mean millions of dollars from more program features or improved utilization."

• "In order to deploy software at Google scale, engineers will touch a good dozen of programming and configuration languages, all with their own interesting optimization problems (from a compiler writer's point of view). Fundamental knowledge in language, compiler, and runtime implementation will help make better engineering decisions everywhere."

• "Did you know that many of our first today's most celebrated engineers have compiler backgrounds? Jeff Dean, Sanjay Ghemawat, Urs Hoelzle, and many others. It's not a coincidence. Compiler optimization trains in big-fiddling as well as algorithmic thinking, which are essential to success in today's world."
The Importance of Compilers

We have hundreds of different programming languages. How do you think they are executed?

And we have hundreds of different hardware. How can we talk to all these different worlds?

But… can you think on why it would do good to you, from a personal point of view, to learn compilers?
Why to Learn Compilers

1. We can become better programmers once we know more about compilation technology.

2. There are plenty of very good job opportunities in the field of compilers.
   - "We do not need that many compiler guys, but those that we need, we need them badly." ♦

3. We can become better computer scientists if we have the opportunity to learn compilation technology.
   - A microcosm crowded with different computer science subjects.
   - Lots of new things happening all the time.

♦: quote attributed to François Bodin, professor at the University of Rennes 1. Previously senior engineer at CAPS
Compilers Help us to be Better Programmers

Compilers usually have different optimization options. The iconic gcc –O1, for instance, runs these optimizations.

What does each of these things do?

We can even run these optimizations individually:

$> gcc -fdefer-pop -o test test.c

We can enable a few of them, and disable others, e.g.:

$> gcc -O1 -fno-defer-pop -o test test.c
Knowledge fixes Misconceptions

“I am a conscious programmer, because I am reusing the name 'i'. In this way my program will consume less memory.”

```
int i = read();
if (i != EOF)
    i = read();
printf("%d", i);
```

“I will not use inheritance in Java, because it can make my method calls more expensive, as I have to find them in the class hierarchy.”

“I will use macros, instead of functions, to have a more efficient code. In this way, there will be no time lost in function calls.”

```
#define MAX(X, Y) (X) > (Y) ? (X) : (Y)
int max(int i, int j) { return i > j ? i : j; }
```
Lots of Job Opportunities

- Expert compiler writers find jobs in many large companies, mostly in the US, Canada, Europe and Japan: Oracle, NVIDIA, Microsoft, Sony, Apple, Cray, Intel, Google, Coverity, MathWorks, IBM, AMD, Mozilla, etc.
- These jobs usually ask for C/C++ expertise.
- Good knowledge of basic computer science.
- Advanced programming skills.
- Knowledge of compiler theory is a big plus!
- Papers published in the field will help a lot too!

Why C/C++ are so important among compiler writers?

compilers.org

Sponsored by Nullstone Corporation - Developers of the
NULLSTONE Automated Compiler Performance Analysis Suite

Compiler jobs for compiler developers who design and develop parsers, optimizers, codegenerators, assemblers, linkers, debuggers, interpreters, IDE's, and related technologies.
Compiler Writers in the Big Companies

Hardware companies need compilation technology.

Intel => icc
Mozilla => JaegerMonkey, IonMonkey, TraceMonkey
Apple => LLVM
NVIDIA => nvcc
Google => Delvik, V8
Microsoft => visual studio, .NET VM
STMicroelectronics => open 64

Can you think about other examples?
Companies that Sell Compilation Technology

There are companies that sell compilation technology to be used in several ways, e.g., to create a new back-end, to parse big data, to analyze programs for security vulnerabilities, etc.

Coverity is a software vendor which develops testing solutions, including static code analysis tools, for C, C++, Java and C#, used to find defects and security vulnerabilities in source code.

The Associated Compiler Experts (ACE) have developed compilers for over 100 industrial systems, ranging from 8-bit microcontrollers to CISC, RISC, DSP and 256-bit VLIW processor architectures.

PathScale Inc. is a company that develops an optimizing compiler for the x86-64 architecture.

The Portland Group, Inc. (PGI) is a company that produces a set of commercially available Fortran, C and C++ compilers for high-performance computing systems.

Green Hills produces compilers for C, C++, Fortran, and Ada. The compilers target 32- and 64-bit platforms, including ARC, ARM, Blackfin, ColdFire

Can you think about other examples?
Hi Prof Pereira,

Long time don't see. I trust you are well.

Here at Apple, we are investing heavily in our compiler technologies. I was wondering if you might have recommendations for students who are interested in pursuing a career in industrial compiler work. I'd be interested in talking to them about opportunities here.

Regards,

Evan Cheng
Sr Manager, Compiler Technologies, Apple (March'13)
Dear Fernando,

I have read your paper in JIT specialization in JavaScript and particularly your work in the context of IonMonkey has caught my eye. I am working as a Research Scientist at Intel Labs on the River Trail project. We are collaborating with Mozilla on a Firefox implementation. Given your recent work on IonMonkey, maybe one of your students would be interested to join us in Santa Clara to work on this topic?

Regards,

Stephan (May'13)
Compiler Knowledge is Fun (and Awesome)

• Compilers are very complex systems. It really takes a lot of programming skills to write them.
  – Large, robust system, usually coded in a high-performance language, such as C or C++, with lots of interactions with the operating system and the hardware.

• Compilers are backed up by a lot of computer science theory. It is not only programming.
  – Type systems, parsing theory, graph theory, algorithms, algebra, fixed point computations, etc

"The first reason Compiler Construction is such an important CS course is that it brings together, in a very concrete way, almost everything you learned before you took the course." Steve Yegge

*: Steve Yegge has worked at Amazon and Google, among others.
Compilers – A Microcosm of Computer Science

- **Algorithms**: graphs everywhere, union-find, dynamic programming
- **Automata Theory**: DFAs for scanning, parser generators, context free grammars.
- **Algebra**: lattices, fixed point theory, Galois Connections, Type Systems
- **Architecture**: pipeline management, memory hierarchy, instruction sets
- **Optimization**: operational research, load balancing, packing, scheduling

◊: Shamelessly taken from slides by *Krishna Nandivada*, professor at the Indian Institute of Technology ([http://www.cse.iitm.ac.in/~krishna/cs3300/lecture1.pdf](http://www.cse.iitm.ac.in/~krishna/cs3300/lecture1.pdf))
The Allies of the Compiler Writer
Static And Dynamic Analysis

• Compilers have two ways to understand programs:
  – Static Analysis
  – Dynamic Analysis

• Static analyses try to discover information about a program without running it.

• Dynamic analyses run the program, and collect information about the events that took place at runtime.

1) Can you give examples of dynamic analyses?
2) And can you give examples of static approaches?
3) What are the pros and cons of each approach?
Dynamic Analyses

• Dynamic analyses involve executing the program.
  – **Profiling**: we execute the program, and log the events that happened at runtime. Example: `gprof`.
  – **Test generation**: we try to generate tests that cover most of the program code, or that produce some event. Example: `Klee`.
  – **Emulation**: we execute the program in a virtual machine, that takes care of collecting and analyzing data. Example: `valgrind`.
  – **Instrumentation**: we augment the program with a meta-program, that monitors its behavior. Example: `AddressSanitizer`. 
Static Analyses

• In this course we will focus on static analyses.
• There are three main families of static analyses that we will be using:
  – **Dataflow analyses**: we propagate information based on the dependences between program elements, which are given by the syntax of the program.
  – **Constraint-Based analyses**: we derive constraints from the program. Relations between these constraints are not determined explicitly by the program syntax.
  – **Type analyses**: we propagate information as type annotations. This information lets us prove properties about the program, such as progress and preservation.
Dataflow Analyses

- Dataflow analyses discover facts about a program propagating information along the control flow graph of this program.

- Many classic analyses fit into this category:
  - Liveness analyses, reaching definitions, constant propagation, available expressions, very busy expressions, etc.

Any idea on what is a Control Flow Graph?
Constraint-Based Analyses

- Constraint-Based analyses are used when the directions in which we must propagate information cannot be easily discovered from the program's syntax.
- The two most important members of this family are control flow analysis, and pointer analysis.

Can you guess what pointer analysis is good for?
Type Systems

- Types lets us associate information with names used in the program.
  - Types provide a good notation for the specification of analyses.
  - They simplify proofs of correctness.
  - May be solved via unification, or data-flow analyses.

Have you heard of progress and preservation?
The Broad Theory

- The algorithms used in compiler optimization include many different techniques of computer science
  - Graphs are everywhere
  - Lattices and the Fixed point theory
  - Many different types of induction
  - Dynamic programming techniques
  - Type theory
  - Integer Linear Programming
  - etc, etc, etc

And all of this in industrial strength compilers!
Graphs

- Graphs are the core of computer science, and permeate code optimization theory.
  - Control Flow Graphs
  - Constraint Graphs
  - Dependence Graphs
  - Strongly Connected Components
  - Graph Coloring

What are each of these graphs? Where are they used?
Fixed Point Theory

- Most of the algorithms that we see in code optimization are iterative. How can we prove that they terminate?
- If we always obtain more information after each iteration of our algorithm...
- and the total amount of information is finite...
- Then, eventually our algorithm must stabilize.
Induction all the way

• Most of our proofs are based on induction.
• We use three main types of induction:
  – Structural induction
  – Induction on derivation rules
  – Induction on size of syntactic terms

The main advantage of proofs by induction is that we already have technology to verify if these proofs are correct mechanically.
The Program Representation Zoo

• Depending on how we represent a program, we may facilitate different static analyses.

• Many dataflow analyses can be solved more efficiently in Static Single Assignment (SSA) form programs.
  – That is why this program representation is nowadays used in almost every compiler that is mildly important.
  – As an example, the tainted flow problem used to be solved by cubic algorithms until the late 2010's, when a faster, quadratic algorithm was proposed, based on a representation called e-SSA form.

• Minimal register assignment in SSA form programs has polynomial time solution, whereas it is NP-complete in general programs!

◇: Tainted Flow Analysis on e-SSA-Form Programs ◇: Trust in the lambda-Calculus
Open Source Community

• Many important compilers are currently open source.
• The Gcc toolkit has been, for many years, the most used compiler for C/C++ programs.
• LLVM is one of the most used compilers for research and in the industry.
• Mozilla's Monkey (SpiderMonkey, TraceMonkey, JagerMonkey, IonMonkey) family of compilers has a large community of users.
• Ocelot is used to optimize PTX, a program representation for graphics processing units.
• The Glasgow Haskell Compiler is widely used in functional programming research.
Conferences and Journals

• There are many important conferences that accept compiler related papers:
  – PLDI: Programming Languages Design and Implementation (10.75, 18%)
  – POPL: Principles of Programming Languages (8.22, 20%)
  – ASPLOS: Architectural Support for Programming Languages and Operating Systems (11.49, 15%)
  – CGO: Code Generation and Optimization (3.56, 31%)
  – CC: Compiler Construction (2.46, 25%)

• And there is a very important journal in this field: TOPLAS – ACM Transactions on Programming Languages and Systems.
A BRIEF HISTORY OF OPTIMIZING COMPILERS
The history of compilers is tightly bound with the history of computer science itself. It is fair to say that without compilation technology the progress of informatics would have been much slower, and probably we would not have internet, www, social networks, gene sequencing, fancy video games, or anything that requires some sort of non-trivial coding.

Can you point some milestone in the history of compilers?
1) What were the first compilers?

2) What were the “first compilation challenges”?

3) Who were the important people in compilation?

4) Which important advances in compilation technology and theory appeared in France?
The Dawn of the First Compilers

• A more serious effort to move the task of generating code away from programmers started in the 50's.

• Fortran was one of the first programming languages, still in use today, to be compiled by an optimizing compiler.

• The developers of Fortran's compiler had to deal with two main problems:
  – Parsing
  – Code optimization
    • register allocation

Fortran was designed to be easily compiled. How so?
Early Code Optimizations

• Frances E. Allen, working alone or jointly with John Cocke, introduced many of the concepts for optimization:
  – Control flow graphs
  – Many dataflow analyses
  – A description of many different program optimizations
  – Interprocedural dataflow analyses
  – Worklist algorithms
• A lot of these inventions and discoveries have been made in the IBM labs.
The Dataflow Monotone Framework

• Most of the compiler theory and technology in use today is based on the notion of the dataflow monotone framework.
  – Propagation of information
  – Iterative algorithms
  – Termination of fixed point computations
  – The meet over all paths solution to dataflow problems

• These ideas came, mostly, from the work of Gary Kildall, who is one of the fathers of the modern theory of code analysis and optimization.

In addition of being paramount to the development of modern compiler theory, Gary Kildall used to host a talk show called "The Computer Chronicles". Nevertheless, he is mostly known for the deal with the Ms-DOS system that involved IBM and Bill Gates.
Abstract Interpretation

• Cousot and Cousot have published one of the most important papers in compiler research\textsuperscript{\textcircled{\textdagger}}, giving origin to the technique that we call Abstract Interpretation.

• Abstract interpretation gives us information about the static behavior of a program.

• We could interpret a program to find out if a property about it is true.
  – But the program may not terminate, and even if it does, this approach could take too long.
  – So, we assign abstract states to the variables, plus an operator called widening, that ensures that our interpretation terminates.
  – This approach may be conservative, but it does terminate!

\textsuperscript{\textcircled{\textdagger}}: Abstract Interpretation: a Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints
Register Allocation

• Register allocation has been, since the early days of compilers, one of the most important code optimizations.

• Register allocation via graph-coloring was introduced by Gregory Chaitin in 1981.

• Linear scan, a register allocation algorithm normally used by JIT compilers, was introduced by Poletto and Sarkar in 1999.

• Linear scan and/or graph coloring are present in most of the modern compilers.
Static Single Assignment

• Once in a while we see smart ideas. Perhaps, the smartest idea in compiler optimization was the Static Single Assignment Form.
  – A program representation in which every variable has only one definition site.
• SSA form was introduced by Cytron et al., in the late eighties in IBM.
• There were many improvements since then, such as pruned SSA form, SSA-based register allocation, etc.
• The idea took off very quickly. Today almost every compiler uses this intermediate representation.

(SSA Seminar): celebrated the 20th anniversary of the Static Single Assignment form, April 27-30, Autrans, France

♡: An Efficient Method of Computing Static Single Assignment Form
Constraint Based Analyses

• Control flow analysis was introduced by Olin Shivers in PLDI'88.
• Pointer analysis is the offspring of many fathers.
  – In 1996, Bjarne Steensgaard described a less precise pointer analysis that could be solved in almost linear time.
  – There is still much research in points-to analyses.

1) Which problems do we solve with control flow analysis? E.g., 0-CFA?
2) What is pointer analysis?
3) Which problems do we solve with pointer analyses?
Type Theory

• Types have been around for a long time.
  – Philosophers and logicians would rely on types to solve paradoxes in Set Theory.

• A lot of work has been done by researchers in the functional programming field, such as Philip Wadler.

• Benjamin Pierce wrote a book, "Types and Programming Languages", in the early 2000's, that has been very influential in the field.

• A major boost in the mechanical verification of theorems is due to several independent groups, such as Xavier Leroy’s (the CompCert compiler) and Frank Pfenning's (the Twelf system).
THE FUTURE

"KEEP LOOKING UP ...
THAT'S THE SECRET OF
LIFE ... "

Snoopy
The Ever Increasing Gap

The Programming Languages are always evolving, usually towards higher levels of abstraction. And the hardware is also always evolving, towards greater performance.

The compiler must bridge this ever increasing gap.

1) How the languages of the future will look like?

2) How the hardware of the future will look like?

3) Where is the research in compilers heading?

4) Is the compiler expert becoming more or less important?
The Challenges of Today

• Parallelism
• Dynamic Languages
• Correctness
• Security

1) Why is this problem an issue?
2) Who is trying to solve this problem?
3) What are the current solutions?
4) Which conference accepts results in that field?
5) What are the unsolved questions?
The Future of Compiler Optimization

• The Full-Employment Theorem ensures that compiler writers will have a lot of work to do in the days ahead.

• Some influent researchers believe that research and development in the field will be directed towards two important paths, in the coming years:
  – Automatic parallelization of programs.
  – Automatic detection of bugs.

• Given that everything happens so fast in computer science, we are likely to see these new achievements coming!
  – And they promise to be super fun 😊

©: Hall, Padua e Pingali “Compiler Research: The Next Fifty Years” Communications of the ACM 2009
The Complete Text (by Snoopy)

It Was A Dark And Stormy Night

Part I
It was a dark and stormy night. Suddenly, a shot rang out! A door slammed. The maid screamed.

Suddenly, a pirate ship appeared on the horizon!

While millions of people were starving, the king lived in luxury. Meanwhile, on a small farm in Kansas, a boy was growing up.

Part II

A light snow was falling, and the little girl with the tattered shawl had not sold a violet all day.

At that very moment, a young intern at City Hospital was making an important discovery. The mysterious patient in Room 213 had finally awakened. She moaned softly.

Could it be that she was the sister of the boy in Kansas who loved the girl with the tattered shawl who was the daughter of the maid who had escaped from the pirates?

The intern frowned.

"Stampede!" the foreman shouted, and forty thousand head of cattle thundered down on the tiny camp. The two men rolled on the ground grappling beneath the murderous hooves. A left and a right. A left. Another left and right. An uppercut to the jaw. The fight was over. And so the ranch was saved.

The young intern sat by himself in one corner of the coffee shop. He had learned about medicine, but more importantly, he had learned something about life.

THE END