CAP - Exercises: static semantics (chapter 4: typing)

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Oct. 2016

Abstract syntax

Recall the abstract syntax of the course for expressions :

$$\begin{array}{rcll} e & ::= & c & \textit{constant} \\ & | & x & \textit{variable} \\ & | & e + e & \textit{addition} \\ & | & e \times e & \textit{multiplication} \\ & | & \ldots & \end{array}$$

and the mini-while language :

$$\begin{array}{rcll} S(Smt) & ::= & x := expr & \text{assign} \\ & | & skip & \text{do nothing} \\ & | & S_1; S_2 & \text{sequence} \\ & | & \texttt{if } b \texttt{ then } S_1 \texttt{ else } S_2 & \text{test} \\ & | & \texttt{while } b \texttt{ do } S \texttt{ done} & \text{loop} \end{array}$$

We expand the language with variable declarations :

$$D(decl) \quad ::= \quad \texttt{var } x : t \quad \text{type declaration}$$

We recall that environnements associate a type to variables (Γ). Here, the environnement is constructed by the following rules :

Declarations

$$\frac{}{\texttt{var } x : t \rightarrow_d [x \mapsto t]}$$

$$\frac{D_1 \rightarrow_d \Gamma_1 \quad D_2 \rightarrow_d \Gamma_2 \quad Dom(\Gamma_1) \cap Dom(\Gamma_2) = \emptyset}{D1; D_2 \rightarrow_d \Gamma_1 \cup \Gamma_2}$$

Expressions Like in the course, for instance :

$$\frac{\Gamma \vdash e_1 : \texttt{int} \quad \Gamma \vdash e_2 : \texttt{int}}{\Gamma \vdash e_1 + e_2 : \texttt{int}}$$

Commands Like in the course, for instance :

$$\frac{\Gamma \vdash b : \texttt{boolean} \quad \Gamma \vdash S : \texttt{void}}{\Gamma \vdash \texttt{while } b \texttt{ do } S \texttt{ done} : \texttt{void}}$$

And a program

$$\frac{D \to \Gamma \quad \Gamma \vdash C : \texttt{void}}{\Gamma \vdash DC : \texttt{void}}$$

EXERCISE ▶ **Well typed**
Type the program :

```
var x1 : integer ; var x2 : integer ; var x3 : integer
x1 := 3 ;
while (not x3) do
x1 := x2 + 1 ;
x3 := x3 and true
done
```

EXERCISE ▶ **Expand expressions**

Complete the abstract syntax and the static semantics (typing) of expressions with the new construction `e1? e2: e3` : if e_1 is true then the expression has value e_2 else e_3.

EXERCISE ▶ **Expand the statements**

Complete the abstract syntax and the static semantics (typing) of statements with an extended for :

```
for i in e1 .. e2 S
```

2 cases :

— The instruction declares the i variable (like in Ada)

— i should be declared before.