

CAP - Exercises: (3-address) Code Generation (chapter 5)

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1 Rules for three address code generation

The code we generate will have an unbounded number of temporaries ($\text{tmp}0, \text{tmp}1, \dots$) but actual LC-3 instructions (ADD, AND, BR ...).

The code generation functions have the following signatures (pseudo-code is given in Figures 1 and 2.) :

$\text{GenCodeExpr} : \text{AExp} \rightarrow \text{Code}^* \times \mathbb{N}$

$\text{GenCodeSmt} : \text{Inst} \rightarrow \text{Code}^*$

where Code^* is a sequence of 3-address instructions (LC-3 with temporaries). As a side effect, the code generation for statements might update a map $\text{Var} \rightarrow \mathbb{N}$ (program variable to a temporary where to find its current value).

Auxiliary functions :

$\text{newTemp}() : \rightarrow \mathbb{N}$

$\text{newLabel}() : \rightarrow \mathbb{N}$

2 Exercises

EXERCISE ▶ By hand !

Using the code generation rules, generate the three-address code for the following (mini-while) program :

```
x1 = 3;  
x2 = 7 + x1;  
while not (x2 < x1) do  
    x2 = x2-1;
```

3 Language expansions

EXERCISE ▶ Rules for boolean expressions

Write a code generation rule for the xor boolean operator.

EXERCISE ▶ Rules for statements

Write a code generation rule for the repeat S until e statement.

e ::= c	<pre>#not valid if c is too big dr <-newTemp() code.add(InstructionAND(dr, dr, 0)) code.add(InstructionADD(dr, dr, c)) return dr</pre>
e ::= x	<pre>#get the place associated to x. regval=getTemp(x) return regval</pre>
e ::= $e_1 + e_2$	<pre>t1 <- GenCodeExpr(e_1) t2 <- GenCodeExpr(e_2) dr <- newTemp() code.add(InstructionADD(dr, t1, t2)) return dr</pre>
e ::= $e_1 - e_2$	<pre>t1 <- GenCodeExpr(e_1) t2 <- GenCodeExpr(e_2) dr <- newTemp() code.add(InstructionNOT(dr, t2)) code.add(InstructionADD(dr, dr, 1)) code.add(InstructionADD(dr, dr, t1)) return dr</pre>
e ::= true	<pre>dr <-newTemp() code.add(InstructionAND(dr, dr, 0)) code.add(InstructionADD(dr, dr, 1)) return dr</pre>
e ::= $e_1 < e_2$	<pre>dr <-newTemp() t1 <- GenCodeExpr (e1-e2) #last write in register (lfalse,lend) <- newLabels() code.add(InstructionBRzp(lfalse)) #if =0 or >0 jump! code.add(InstructionAND(dr, dr, 0)) code.add(InstructionAND(dr, dr, 1)) #dr <- true code.add(InstructionBR(lend)) code.addLabel(lfalse) code.add(InstructionAND(dr, dr, 0)) #dr <- false code.addLabel(lend) return dr</pre>

FIGURE 1 – Code generation for expressions

x := e	<pre> dr <- GenCodeExpr(e) #a code to compute e has been generated if x has a location loc: code.add(instructionADD(loc,dr,0)) else: storeLocation(x,dr) </pre>
S1 ; S2	<pre> #concat codes GenCodeSmt(S1) GenCodeSmt(S2) </pre>
if b then S1 else S2	<pre> dr <-GenCodeExpr(b) #dr is the last written register lfalse,lendif=newLabels() code.add(InstructionBRz(lfalse) #if 0 jump to execute S2 GenCodeSmt(S1) #else (execute S1 code.add(InstructionBR(lendif)) #and jump to end) code.addLabel(lfalse) GenCodeSmt(S2) code.addLabel(lendif) </pre>
while b do S done	<p>TODO!!</p>

FIGURE 2 – Code generation for Statements