

# CAP - Exercises: CFG and data-flow analyses (chapter 7)

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## 1 Liveness analysis

A variable at the left-hand side of an assignment is *killed* by the block. A variable that appears in this bloc is *generated*.

$$LV_{exit}(\ell) = \begin{cases} \emptyset & \text{if } \ell = \text{final} \\ \bigcup\{LV_{entry}(\ell') | (\ell', \ell) \in \text{flow}(G)\} & \text{otherwise} \end{cases}$$

$$LV_{entry}(\ell) = (LV_{exit}(\ell) \setminus kill_{LV}(\ell)) \cup gen_{LV}(\ell)$$

### EXERCISE ► Live variables

Generate the CFG for the following program :

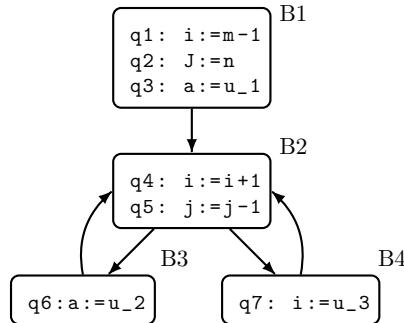
```
while d>0 do {
    a:=b+c;
    d:=d-b;
    e:=a+f;
    if e>0 then {
        f:=a-d;
        b:=d+f;
    }
    else{
        e:=a-c;
    }
    b:=a+c;
}
```

On this CFG :

- Write and solve the equations defining the sets *Gen*, *Kill*, *In* and *Out* for the liveness analysis, pay attention to initialisations.
- Suppress the dead code.

### EXERCISE ► Live Variables

After code generation, we obtain the following graph :



On this graph, perform liveness analysis and suppress the dead code.

## 2 Available expressions

We recall :

$$AE_{entry}(\ell) = \begin{cases} \emptyset & \text{if } \ell = init \\ \cap\{AE_{exit}(\ell') | (\ell', \ell) \in \text{flow}(G)\} & \end{cases}$$

$$AE_{exit}(\ell) = (AE_{entry}(\ell) \setminus kill_{AE}(\ell)) \cup gen_{AE}(\ell)$$

### EXERCISE ▶ Common (sub) expressions

On the following 3-address code :

```

(1)   a=b+c
(2)   b=a+c
      d=c+e
      si d>7 aller a (8)
      t=a+c
      v=c+e
      aller a (10)
(8)   t=b+c
      v=a+c
(10)  b=a+c
      a=b+c

```

1. Construct the CFG.
2. For each block, compute the sets `gen` and `kill` sets for common subexpressions.
3. Compute the set of all available expressions at the entry and exit of each block.
4. Optimise the code.

### EXERCISE ▶ With a loop

Same questions with :

```

x:=a+b;
y:=a*b;
while(y>a+b) do
    a:=a+a;
    x:=a+b;
done

```



