CAP - Exercise: Op. semantics for procedures (chapter 8)

Laure Gonnord

Nov. 2016

Abstract syntax :

```
\begin{array}{rcl} S & \in & \mathbf{Stm} \\ S & ::= & x := a | \mathtt{skip} | S_1; S_2 \mid \\ & & \texttt{if } b \texttt{ then } S_1 \texttt{ else } S_2 \mid \\ & & \texttt{while } b \texttt{ do } S \mid \\ & & \texttt{begin } D_V D_P; S \texttt{ end } \mid \\ & & \texttt{call } p \\ D_V & ::= & \texttt{var } x := a; D_V | \varepsilon \\ D_P & ::= & \texttt{proc } p \texttt{ is } S; D_P | \varepsilon \end{array}
```

<u>EXERCISE</u> \blacktriangleright Operational semantics

For the following miniwhile with procedures program :

```
(0) begin
(1)
        var x:=0
(2)
        proc p is x:=x*2
(3)
        proc q is call p
(4)
        begin
(5)
              var x:=5
(6)
              proc p is x:=x+1
(7)
              call q
(8)
              y := x
(9)
        end
(10) end
```

Apply the semantic rules of the course with :

— A Dynamic binding for variables and procedures.

— Dynamic binding for variables, and static for procedures.

$\underline{\text{EXERCISE}} \triangleright \mathbf{Procedures with parameters}$

Add a new abstract syntax for the mini-while with procedures with two parameters. Give a formal "call by value" semantics (environment, update, call). Give an example.